

**Exam – Sample Questions**  
**CS 151: Algorithms and Data Structures**

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NOTE: the actual exam will contain around 20-25 questions.

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1. Consider a LinkedList that has a `get(int i)` method to return the  $i$ -th element in the list. Write pseudocode to implement the method. Next, analyse (in order-notation, in terms of  $n$ ) the following code:

```
LinkedList L;
// ... LinkedList is initialized and now contains n items (code not shown).
for (int i=1; i<=n; i++) {
    sum = sum + L.get(i);
}
```

2. Consider this snippet of code:

```
int[][] A = new int [n][];
for (int i=0; i<n; i++) {
    if (i % 2 == 0) // i is a multiple of 2
        A[i] = new int [n];
    else
        A[i] = new int [1];
}
```

```
for (int i=0; i<A.length; i++)
    for (int j=0; j<A[i].length; j++)
        sum = sum + A[i][j];
```

It's running time is:

- (a)  $O(n \log n)$ .
  - (b)  $O(n)$ .
  - (c)  $O(n^2)$ .
  - (d)  $O(\frac{n}{2})$ .
3. Suppose  $f(n) = 3n^3 + 2n^2 + n$ . Consider these statements:
- I  $f(n) = O(n^3)$ .
  - II  $f(n) = O(n^4)$ .

- (a) Only I is true.
  - (b) Only II is true.
  - (c) I and II are both true.
  - (d) I and II are both false.
4. The ratio of the running time of Bubble-Sort to that of Merge-Sort is:
- (a)  $O(n \log n)$ .
  - (b)  $O(\frac{n}{\log n})$ .
  - (c)  $O(\log \frac{n^2}{\log n})$ .
  - (d)  $O(1)$ .

5. Consider the following input with 10 keys (each is a letter) to QuickSort.

K G I K N V S S W Q

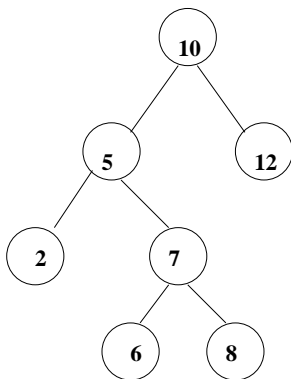
Assume the leftmost element is the partitioning element. Show the array after the partitioning step is complete.

6. Suppose the indices of the array that implements a binary heap start from 1. Now consider these assertions about finding the indices of the parent and left-child of the  $i$ -th element:

- I parent( $i$ ) =  $i/2$ .
- II left( $i$ ) =  $2i + 1$ .

- (a) I and II are both true.
- (b) Only I is true.
- (c) Only II is true.
- (d) I and II are both false.

7. Delete 5 from this binary tree:



8. Insert the bitstrings 100, 110, 011 and 010 into a simple trie.
9. Consider the statements
- I In a simple trie, all the keys are at the leaves.
  - II In a full trie, all the keys are at the leaves.
- (a) I and II are both true.
  - (b) Only I is true.
  - (c) Only II is true.
  - (d) I and II are both false.
10. Draw an NFA for the regular expression  $C(AB)^*(A|C)$
11. Suppose the height of a binary tree with one node is defined to be 1. A complete binary tree of height  $h$  has
- (a)  $2^h$  edges.
  - (b)  $2^h - 1$  edges.
  - (c)  $2^h - 2$  edges.
  - (d)  $2^{h+2}$  edges.
12. A ring is a connected undirected graph in which all the nodes form a ring (thus, each node has exactly two neighbors). The running time of Prim's algorithm, using an adjacency list, on this graph is:
- (a)  $O(V \log V)$
  - (b)  $O(V^2 \log V)$
  - (c)  $O(V^2)$
  - (d)  $O(V^3)$
13. Consider these two statements about a connected undirected graph with  $V$  vertices and  $E$  edges:
- I  $O(V) = O(E)$
  - II  $O(E) = O(V^2)$
- (a) I and II are both false.
  - (b) Only I is true.
  - (c) Only II is true.
  - (d) I and II are both true.

14. Consider the following two statements about an undirected graph with  $n$  vertices:
- I If every pair of vertices has an edge, there are exactly  $\frac{n(n+1)}{2}$  edges.
  - II If the graph has a spanning tree, the tree must have  $\Theta(n)$  edges.
- (a) I and II are both true.
  - (b) Only I is true.
  - (c) Only II is true.
  - (d) I and II are both false.
15. Suppose the shortest path from node  $i$  to node  $j$  goes through node  $k$  and that the cost of the subpath from  $i$  to  $k$  is  $D_{ik}$ . Consider these two statements:
- I Every shortest path from  $i$  to  $j$  must go through  $k$ .
  - II Every shortest path from  $i$  to  $k$  has cost  $D_{ik}$ .
- (a) I and II are both true.
  - (b) Only I is true.
  - (c) Only II is true.
  - (d) I and II are both false.
16. Consider the contiguous load balancing problem where  $s_i$  is the execution time of task  $i$  and let  $D_i^k$  be the optimal cost in assigning tasks  $0, \dots, i$  on to  $k$  processors. Then, the dynamic programming recurrence can be stated as:
- (a)  $D_i^k = \max_j \min(D_j^{k-1}, \sum_{l=j+1}^i s_l)$ .
  - (b)  $D_i^k = \max_j \min(D_{j-1}^k, \sum_{l=j+1}^i s_l)$ .
  - (c)  $D_i^k = \max_j \max(D_j^{k-1}, \sum_{l=j+1}^i s_l)$ .
  - (d)  $D_i^k = \min_j \max(D_j^{k-1}, \sum_{l=j+1}^i s_l)$ .
17. Consider the execution times of two algorithms I and II:
- I  $O(n^{\log n})$
  - II  $O(\log(n^n))$
- (a) Only I is polynomial.
  - (b) Only II is polynomial.
  - (c) I and II are both polynomial.
  - (d) Neither I nor II is polynomial.
18. For a travelling salesman problem with  $n$  cities, the number of possible tours is:

- (a)  $\frac{n!}{2}$
- (b)  $\frac{(n-1)!}{2}$
- (c)  $\frac{n(n-1)}{2}$
- (d)  $\frac{n(n+1)}{2}$

19. Consider the following statements about NP-completeness:

- I NP-completeness only applies to combinatorial optimization problems.
- II An NP-complete problem cannot be solved in polynomial-time.

- (a) Only I is true.
- (b) Only II is true.
- (c) I and II are both true.
- (d) Neither I nor II is true.

20. Consider the following statements about NP-completeness:

- I The traveling salesman problem is NP-complete.
- II The minimum-spanning tree problem is not NP-complete.

- (a) Only I is true.
- (b) Only II is true.
- (c) I and II are both true.
- (d) Neither I nor II is true.

21. Consider the following statements about estimation:

- I The sample mean is always larger than the sample variance.
- II It is possible for the sample variance to be zero.

- (a) Only I is true.
- (b) Only II is true.
- (c) I and II are both true.
- (d) Neither I nor II is true.